

# CS151 Intro to Data Structures

Queues

Lists

# Announcements

- HW02 due last night
  - two days remaining for late submissions (10 points off each day)
  - you cannot assume the csv files are in the same directory
    - <https://piazza.com/class/lr8bl2y6e4f5ql/post/55>
      - `java DriverHW02 -f Dianna ../../names/names1990.csv`
    - other issues? ask me during lab
    - <https://www.menti.com/ale73dwub17g>
    - 2918 6510
  - HW03 (Stacks & Queues) released – due Thursday 2/22
    - easier than hw2

# Outline

- Junit review
- Stacks & Queues review
- Dequeues
- ADTs
- Lists
- Iterators

# Junit

# Using JUnit

Import Test Annotation Framework

```
import org.junit.Test;
```

- Write tests using `@Test` annotation

```
@Test
public void testEmpty() {
    ArrayStack<String> stack = new ArrayStack<String>(10);
    assertTrue(stack.isEmpty());
}
```

- Run with the `run_junit.sh` script in lecture-notes

# Junit

- How to write good tests
  - Try to test every branch “branch coverage”
  - aim to cover every branch, but this is hard! ~90% is ok

```
1. public void insert(E item, int index) {
2.     if (index > numElems) {
3.         throw new ArrayIndexOutOfBoundsException();
4.     }
5.
6.     if (numElems >= data.length) {
7.         expand();
8.     }
9.
10.    ...
11. }
```

# Stacks Review - FILO

- First In Last Out
- *stack* of plates in the dining hall
  
- Big O of the following operations?
  - push
    - **depends on how we implemented!**
    - LL?  $O(1)$
    - ExpandableArray?  $O(n)$  worst case
  - pop
    - LL?  $O(1)$
    - ExpandableArray?  $O(1)$
  - peek
    - LL?  $O(1)$
    - ExpandableArray?  $O(1)$
  - isEmpty?
    - LL?  $O(1)$
    - ExpandableArray?  $O(1)$

# Queues - First-in First-out

The first item in, is the first item out (grocery line)



# Queue Interface

```
public interface Queue<E> {  
    int size();  
    boolean isEmpty();  
    E first();  
    void enqueue(E e);  
    E dequeue();  
}
```

- **null is returned from dequeue() and first() when queue is empty**

# Implementing a Queue with an Array

- enqueue?  $O(1)$
- dequeue?  $O(n)$ 
  - How can we fix this with an easy trick?
  - Hint: we have a lastElemIdx... what other field can we add?
- isEmpty?  $O(1)$
- first?  $O(1)$

# Implementing a Queue with a LinkedList

- enqueue?
  - runtime complexity?
    - $O(1)$
- dequeue?
  - runtime complexity?
    - $O(1)$
- first?
  - runtime complexity?
    - $O(1)$
- isEmpty?
  - runtime complexity?
    - $O(1)$

# Performance and Limitations

- Performance
  - let  $n$  be the number of objects in the queue
  - The space used is  $O(n)$
  - Each operation runs in time  $O(1)$
- Limitations
  - Depending on the implementation,
    - max size is limited and can not be changed
    - Or need to grow the array when out of room

Deque / “Decks”

# Queues (FIFO)

- Restaurant Waitlist
  - FIFO make sense!
  - What if we pop off a party in the waitlist but then a table wasn't ready... we need to add them back in... push will put them at the end!
- Solution? **Deque**
  - gives us more flexibility

```
public interface Deque<E> {

    //returns number of elements in the dequeu
    int size();

    //returns true if the dequeu is empty, false otherwise
    boolean isEmpty();

    //returns top element in the dequeu (or null if empty)
    E first();

    //returns top element in the dequeu (or null if empty)
    E last();

    //inserts the element e to the beginning
    void addFirst(E e);

    //adds the element e to the end
    void addLast(E e);

    //returns the first element in the dequeu (or null if empty)
    E removeFirst();

    //returns the last element in the dequeu (or null if empty)
    E removeLast();
}
```

# Dequeues

- Implemented as array with:
  - front marker
  - rear marker



# Front and Back Markers

addFirst:

$$f = (f - 1 + n) \% n;$$

addLast:

$$r = (r + 1) \% n;$$

removeFirst:

$$f = (f + 1) \% n;$$

removeLast:

$$r = (r - 1 + n) \% n;$$

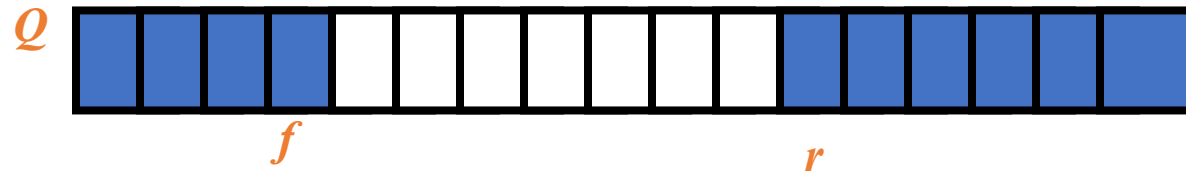
where  $n$  is size of the array

# Circular Queue

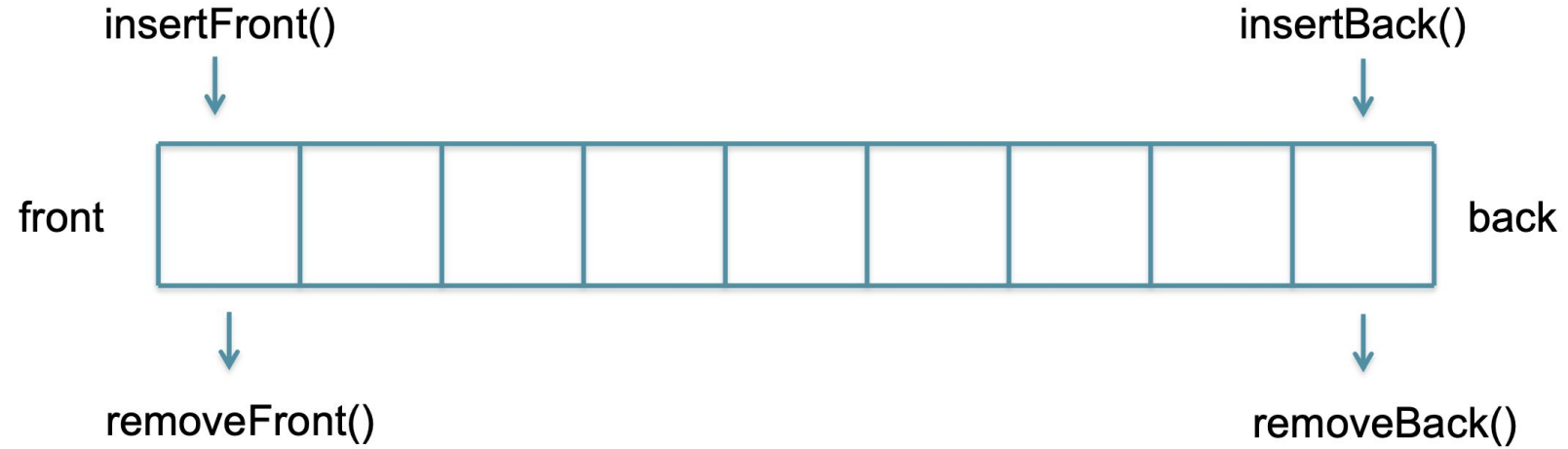
$r = (r+1) \% n$  is the first empty slot past the rear of the queue

where  $n = \text{length of the array}$

wrapped-around configuration



# Doubled-ended Queue (aka Deques aka “Decks”)



## ***Dynamic Data Structure used for storing sequences of data***

- Insert/Remove at either end in  $O(1)$
- If you exclusively add/remove at one end,
  - then ***it becomes a stack***
- If you exclusive add to one end and remove from other,
  - then ***it becomes a queue***

ADTs

# Abstract Data Types

- high-level description of a set of operations that can be performed on a data structure
- It defines the behavior of a data type independently of its implementation
- Cannot instantiate
- What does this remind you of that we've learned so far?

# Queue ADT

<https://docs.oracle.com/javase/8/docs/api/java/util/Queue.html>

Look at the “Implementing Classes”

# Abstract Data Types

- There are multiple ways to implement a data structure each with different trade offs
  - Ex. stack can be implemented with an array or a linked list
- **List ADT:**
  - supports a *linear sequence of elements*

# Lists



# java.util.List ADT

- `size()`: Returns the number of elements in the list.
- `isEmpty()`: Returns a boolean indicating whether the list is empty.
- `get(i)`: Returns the element of the list having index *i*; an error condition occurs if *i* is not in range  $[0, \text{size}() - 1]$ .
- `set(i, e)`: Replaces the element at index *i* with *e*, and returns the old element that was replaced; an error condition occurs if *i* is not in range  $[0, \text{size}() - 1]$ .
- `add(i, e)`: Inserts a new element *e* into the list so that it has index *i*, moving all subsequent elements one index later in the list; an error condition occurs if *i* is not in range  $[0, \text{size}()]$ .
- `remove(i)`: Removes and returns the element at index *i*, moving all subsequent elements one index earlier in the list; an error condition occurs if *i* is not in range  $[0, \text{size}() - 1]$ .

# Example

Method	Return Value	List Contents
add(0, A)		

# Example

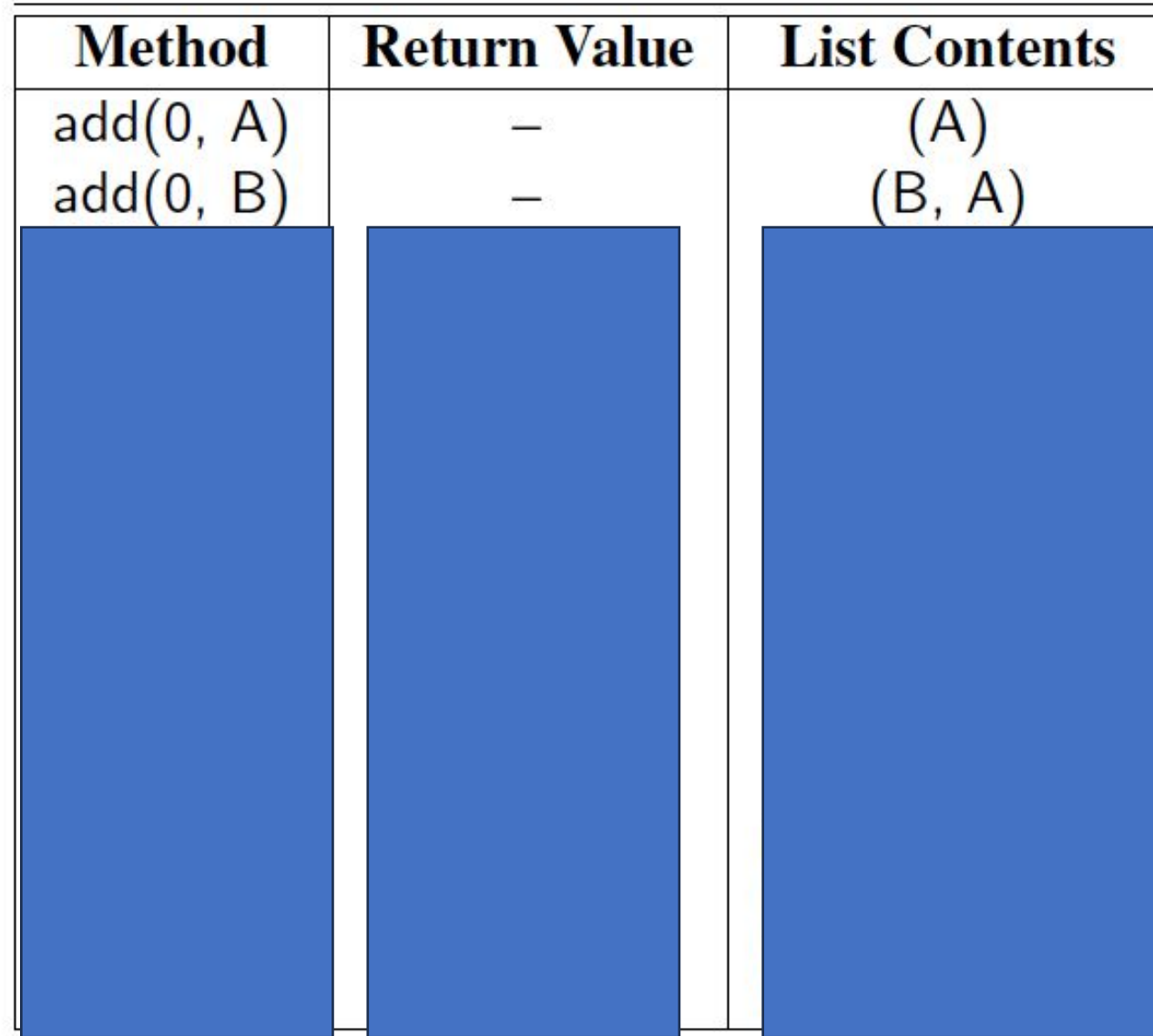
Method	Return Value	List Contents
add(0, A)	–	(A)

# Example

Method	Return Value	List Contents
add(0, A) add(0, B)	—	(A)

# Example

Method	Return Value	List Contents
add(0, A)	–	(A)
add(0, B)	–	(B, A)



# Example

Method	Return Value	List Contents
add(0, A)	–	(A)
add(0, B)	–	(B, A)
get(1)		
set(2, C)		
add(2, C)		
add(4, D)		
remove(1)		
add(1, D)		
add(1, E)		
get(4)		
add(4, F)		
set(2, G)		
get(2)		

# Example

Method	Return Value	List Contents
add(0, A)	–	(A)
add(0, B)	–	(B, A)
get(1)	A	(B, A)
set(2, C)	“error”	(B, A)
add(2, C)	–	(B, A, C)
add(4, D)	“error”	(B, A, C)
remove(1)	A	(B, C)
add(1, D)	–	(B, D, C)
add(1, E)	–	(B, E, D, C)
get(4)	“error”	(B, E, D, C)
add(4, F)	–	(B, E, D, C, F)
set(2, G)	D	(B, E, G, C, F)
get(2)	G	(B, E, G, C, F)

# List ADT

<https://docs.oracle.com/javase/8/docs/api/java/util/List.html>

Look at the “all known implementing classes”

We’re going to focus on ArrayList today



# List ADT

## Reminder of our methods:

- `size()`: Returns the number of elements in the list.
- `isEmpty()`: Returns a boolean indicating whether the list is empty.
- `get(i)`: Returns the element of the list having index *i*; an error condition occurs if *i* is not in range  $[0, \text{size}() - 1]$ .
- `set(i, e)`: Replaces the element at index *i* with *e*, and returns the old element that was replaced; an error condition occurs if *i* is not in range  $[0, \text{size}() - 1]$ .
- `add(i, e)`: Inserts a new element *e* into the list so that it has index *i*, moving all subsequent elements one index later in the list; an error condition occurs if *i* is not in range  $[0, \text{size}()]$ .
- `remove(i)`: Removes and returns the element at index *i*, moving all subsequent elements one index earlier in the list; an error condition occurs if *i* is not in range  $[0, \text{size}() - 1]$ .

# ArrayList

Big-O memory?

- $O(n)$

Indexing / random access?

- $O(1)$

Add / remove?

- $O(n)$

# Iterators

# Iterators

- represents a sequence of elements and provides a way to iterate, or traverse, through those elements one at a time

# Iterators

- Abstracts the process of scanning through a sequence of elements (traversal)
- provides a way to iterate, or traverse, through elements one at a time

`hasNext()`: Returns true if there is at least one additional element in the sequence, and false otherwise.

`next()`: Returns the next element in the sequence.

- Combination of these two methods allow a generic traversal structure

```
while (iter.hasNext ()) {  
    iter.next ();  
}
```

# Iterators

- **code**
- Can an iterator go backwards? NO. Only can do `next()`

# Iterable Interface

- What can i use an `iterator` on? Anything that implements the `iterable` interface.
- Each call to `iterator()` returns a new iterator instance, thereby allowing traversals of a collection
- `List` interface extends `Iterable` and `ArrayList` implements `List`

# Iterable Interface

An interface with a single method:

- `iterator()`: returns an iterator of the elements in the collection



# Iteratoror Interface

# Iterator Interface

Another interface that supports iteration

- `boolean hasNext()`
- `E next()`
- `void remove()`
  
- `Scanner` implements `Iterator<String>`
- `ArrayList` inner class `ArrayListIterator` implements `Iterator`

Let's make `ExpandableArray` iterable

# Iterable versus Iterator?

- Iterable

- `java.lang`
- **override** `iterator()`
- Doesn't store the iteration state
- Removing elements during iteration isn't allowed

- Iterator

- `java.util`
- **Override** `hasNext()`, `next()`
- **Optional** `remove()`
- Stores iteration state (list cursor)
- Removing elements during iteration supported